

MORE ON THE DECOMPOSITION OF TREES INTO ISOMORPHIC SUBTREES

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Abstract

Caro and Schonheim [2] gave a necessary condition for a tree T to be the union of pairwise edge-disjoint subtrees, each isomorphic to a given tree G , and showed that this condition is not sufficient in general.

Answering a question raised in [2], we give a simple characterization of those trees G for which this condition is also sufficient for all trees T .

Our notation is similar to that of Alon in [1]. All graphs considered in this paper are finite and undirected. A graph H is said to have a G -decomposition if it is the union of pairwise edge-disjoint subgraphs, each isomorphic to G . We denote this situation by $G|H$.

We denote by $V(G)$ the set of vertices of G and by $E(G)$ the set of edges of G , and put $e(G) = |E(G)|$. If T is a tree, $C(T)$ is the set of all cut points of T . If $u \in C(T)$ and $\{(u, z_i) : 1 \leq i \leq s\}$ is the set of edges incident with u , then $T-u$ is a forest consisting of trees T_1, \dots, T_s where $z_i \in T_i$ for $1 \leq i \leq s$. The branches of T at u are the trees $T_i \cup (u, z_i)$, $1 \leq i \leq s$. If v is a vertex of T , $v \neq u$, let $B(u, v, T)$ denote the unique branch of T at u that contains v . We denote by $\mathcal{B}(u, T)$ the set of all branches of T at u , and put $\mathcal{B}_v(u, T) = \mathcal{B}(u, T) \setminus \{B(u, v, T)\}$. We also define $m(u, T) = \max \{e(B) : B \in \mathcal{B}(u, T)\}$.

Denote by $d_{t,k}(u, T)$ the number of branches B at u such that $e(B) \equiv t \pmod k$; the $(\text{mod } k)$ branching vector of u is the vector $d_k(u, T) = (d_{1,k}(u, T), d_{2,k}(u, T), \dots, d_{k-1,k}(u, T))$. If T has k edges we write $d(u, T)$ instead of $d_k(u, T)$. Notice that in this case $d_{i,k}(u, T)$ is just the number of branches at u having exactly i edges.

Let G be a tree with k edges, and let T be a tree. Denote by $G||T$ the following condition: For every $v \in C(T)$, $d_k(v, T)$ is a linear combination with nonnegative integer coefficients of the vectors $d(u, G)$ ($u \in C(G)$).

In this notation Theorem 2 in [2] gives the following necessary condition for a tree T to have a G -decomposition.

THEOREM A. *If G is a tree and $e(G) > 1$, then for every tree T*
 $G|T + G||T$.

The converse is not true in general, and some simple counterexamples are given in [2]. In this paper we give a simple characterization of those trees G for which the converse implication holds for all trees T .

Another characterization of these trees G was given by Alon in Theorem 1 of [1] (see condition (ii) below). This characterization is not so simple and, as noted by the referee, it does not yield an easy algorithm to check whether a given tree G satisfies the converse of Theorem A. The characterization given here supplies such an algorithm.

Let G be a tree with k edges, $k > 1$. Consider the following four conditions:

- (i) For every tree T , $G||T + G|T$.
- (ii) For every tree T with k edges, $G||T + G|T$ (i.e., $G||T + T$ is isomorphic to G).
- (iii) For every two distinct vertices $u, v \in C(G)$, the following condition $C = C(u, v)$ holds.

Condition $C(u, v)$: If $e(B(u, v, G)) = e(B(v, u, G))$, then there is a bijection $f: B_v(u, G) \rightarrow B_u(v, G)$ such that for every $B \in B_v(u, G)$ B is isomorphic to $f(B)$ by an isomorphism that carries u onto v . (The existence of such a bijection f is clearly equivalent to the existence of an isomorphism $g: UB_v(u, G) \rightarrow UB_u(v, G)$ that maps u onto v .)

- (iv) For every $u, v \in C(G)$

$$m(u, G) - m(v, G) = d(u, G) - d(v, G).$$

(The equality $d(u, G) = d(v, G)$ is clearly equivalent to the existence of a size-preserving bijection $f: B(u, G) \rightarrow B(v, G)$, i.e., a bijection

f that satisfies $e(B) = e(f(B))$ for all $B \in B(u, G)$.)

We shall prove the following theorem.

THEOREM 1. *Let G be a tree with k edges, $k > 1$. Then conditions (i) - (iv) are equivalent.*

Clearly, it is very easy to check if a given tree G satisfies (iv). (In fact, we can devise an algorithm that will do the checking in $O(k)$ steps, assuming a suitable standard presentation of G .)

We note that the equivalence of (ii) and (iv) answers Question 2 raised in [2]. The equivalence of (i) and (ii) is just the assertion of Theorem 1 in [1]. Thus, in order to prove Theorem 1, it remains to show that (ii), (iii) and (iv) are equivalent.

We need two simple lemmas.

LEMMA 1. *If G and T are trees, $e(G) = e(T) = k > 1$ and $G||T$, then for every $v \in C(T)$ there is a $u \in C(G)$ such that $d(v, T) = d(u, G)$.*

Proof. By definition there are $u_1, \dots, u_m \in C(G)$ and nonnegative integers x_1, \dots, x_m such that

$$(1) \quad d(v, T) = \sum_{i=1}^m x_i d(u_i, G).$$

Define $y = (1, 2, \dots, k-1)$. Clearly, the scalar product of y and $d(v, T)$ is the sum of the sizes of all branches in $B(v, T)$, which is k . Similarly, the scalar product of y and $d(u_i, G)$ is k for all i , $1 \leq i \leq m$. Thus (1) implies

$$k = \sum_{i=1}^m x_i \cdot k$$

and the result follows. \square

LEMMA 2. *Let G be a tree and suppose $u, v \in C(G)$, $u \neq v$. Then*

$$(2) \quad m(u, G) = m(v, G)$$

iff

$$(3) \quad e(B(u, v, G)) = e(B(v, u, G)).$$

If (2) (and (3)) hold, then $B(u,v,G)$ is the unique branch at u having $m(u,G)$ edges.

Proof. Clearly $B(u,v,G)$ properly contains every branch in $\mathcal{B}_u(v,G)$, and thus for every $B \in \mathcal{B}(v,G) \setminus \{B(u,v,G)\}$

$$(4) \quad e(B(u,v,G)) > e(B).$$

Similarly, for every $C \in \mathcal{B}(u,G) \setminus \{B(u,v,G)\}$

$$(5) \quad e(B(v,u,G)) > e(C).$$

Suppose (2) holds. If (3) is false, we may assume, without loss of generality, that

$$(6) \quad e(B(u,v,G)) > e(B(v,u,G)).$$

Relations (4) and (6) imply

$$m(v,G) < e(B(u,v,G)) \leq m(u,G),$$

contradicting (2). Thus (2) implies (3). Conversely, if (3) holds, then (4) and (5) imply that

$$m(u,G) = e(B(u,v,G)) = e(B(v,u,G)) = m(v,G),$$

and thus (3) implies (2).

If (3) holds, then (5) implies that $B(u,v,G)$ is the unique branch at u having $m(u,G)$ edges. \square

Proof of Theorem 1. We have to prove that conditions (ii), (iii) and (iv) are equivalent. We split the proof into four steps.

Step I. (iii) \rightarrow (iv).

Assume (iii) holds and suppose $u, v \in C(C)$ satisfy $m(u,C) = m(v,C)$. We have to show that $d(u,G) = d(v,G)$. If $u = v$ this is trivial. Otherwise, $e(B(u,v,G)) = e(B(v,u,G))$ by Lemma 2. By (iii), there is a bijection $f: \mathcal{B}_v(u,G) \rightarrow \mathcal{B}_u(v,G)$ such that B is isomorphic to

$f(B)$ (and certainly $e(B) = e(f(B))$) for all $B \in \mathcal{B}_v(u,G)$. This, and the fact that $e(B(u,v,G)) = e(B(v,u,G))$, imply that $d(u,G) = d(v,G)$, as needed.

Step II. (iv) \rightarrow (iii).

We suppose that (iv) holds and prove (iii). Assume (iii) is false, and choose among all pairs of cut points of G that violate condition C a pair (u,v) for which $e(B(u,v,G)) (= e(B(v,u,G)))$ is maximal. By Lemma 2 $m(u,G) = m(v,G)$ and thus by (iv), $d(u,G) = d(v,G)$. Since $e(B(u,v,G)) = e(B(v,u,G))$, this means that there is a size-preserving bijection $f: \mathcal{B}_v(u,G) \rightarrow \mathcal{B}_u(v,G)$. However, u and v do not satisfy condition C, and thus there is a branch $B \in \mathcal{B}_v(u,G)$ that is not isomorphic to $f(B)$ by an isomorphism that carries u onto v . Let u_1 be the unique neighbor of u in B and let v_1 be the unique neighbor of v in $f(B)$. Clearly $u_1, v_1 \in C(C)$,

$$(7) \quad e(B(u_1, v_1, G)) = e(C) - e(B) + 1 = e(C) - e(f(B)) + 1 \\ = e(C) - e(B(v_1, u_1, G)),$$

and

$$(8) \quad e(B(u_1, v_1, G)) = e(C) - e(B) + 1 > e(C) - e(B) > e(B(u, v, G)).$$

Because of the extremal choice of u and v (see above), inequality (8) implies condition C(u_1, v_1). Because of (7), this means that there is an isomorphism $g: U(B_{u_1}(u_1, G)) \rightarrow U(B_{u_1}(v_1, G))$ that carries u_1 onto v_1 . But this g can clearly be extended to an isomorphism from B to $f(B)$ that carries u onto v . This contradicts the fact that there is not such an isomorphism and thus completes the proof of Step II.

Step III. (ii) \rightarrow (iv)

Assume (ii) is false. We prove that (ii) is false by constructing a tree T with k edges such that $G|T$ but T is not isomorphic to G . Let $u, v \in C(C)$ satisfy $m(u,C) = m(v,C) (= e(B(u,v,C)))$ but $d(u,G) \neq d(v,G)$. Let T be the tree obtained from G by replacing

the branches in $B_v(u, G)$ by isomorphic copies of the branches in $B_u(v, G)$. Thus T consists of two edge-disjoint subtrees T_0, T_1 with a common vertex u : $T_0 = B(u, v, G)$, and T_1 is isomorphic to $UB_u(v, G)$ by an isomorphism r that carries u onto v . Note that the replacement of $UB_v(u, G)$ by T_1 does not affect the vectors $d(x, G)$ for $x \in V(T_0) \setminus \{u\}$.

We first show that $G \parallel T$. Suppose $y \in C(T)$. If $y \in V(T_0) \setminus \{u\}$ then $d(y, T) = d(y, G)$, as noted above. If $y \in V(T_1) \setminus \{u\}$ then $d(y, T) = d(r(y), G)$. Clearly $d(u, T) = d(v, G)$. Thus $G \parallel T$.

Next we show that T is not isomorphic to G , by showing that the number of cut points $y \in C(T)$ for which $d(y, T) = d(v, G)$ is greater than the number of cut points $x \in C(G)$ that satisfy $d(x, G) = d(v, G)$. Indeed, if $x \in V(G) \setminus V(B(u, v, G))$ then clearly $m(x, G) > e(B(u, v, G)) = m(u, G) = m(v, G)$ and thus $d(x, G) \neq d(v, G)$. In addition $d(u, G) \neq d(v, G)$ and thus all the cut points $x \in C(G)$, for which $d(x, G) = d(v, G)$ belong to $V(B(u, v, G)) \setminus \{u\}$. However, for each such x , $d(x, T) = d(x, G) = d(v, G)$ and in addition $d(u, T) = d(v, G)$. This completes Step III.

Step IV. (iv) \rightarrow (ii)

Suppose (iv) holds. Let T be a tree with k edges, and suppose $G \parallel T$. We have to show that T is isomorphic to G . Assume this is false. By Lemma 1, for every $v \in C(T)$ there is a $u \in C(G)$ and a size-preserving bijection $f: B(u, G) \rightarrow B(v, T)$. Since G is not isomorphic to T , there is a branch $B \in B(u, G)$ such that B is not "properly" isomorphic to $f(B)$, i.e., such that there is no isomorphism $g: B \rightarrow f(B)$ that carries u onto v . Let $n(v, u, f)$ denote the smallest possible number of edges of a branch $B \in B(u, G)$ that is not properly isomorphic to $f(B)$. Define $n = \{\min n(v, u, f)\}$, where the minimum is taken over all triples $\langle v, u, f \rangle$ with $v \in C(T)$, $u \in C(G)$, $d(v, T) = d(u, G)$ and $f: B(u, G) \rightarrow B(v, T)$ a size-preserving bijection. Suppose $n = n(v_0, u_0, f_0)$ and let $B \in B(u_0, G)$ be a branch of size n at u_0 that is not properly isomorphic to $C = f_0(B)$. Let u_1 be the unique neighbor of u_0 in B and let v_1 be the

unique neighbour of v_0 in C . Clearly $v_1 \in C(T)$, since otherwise $e(B) = e(C) = 1$, and B would be properly isomorphic to C . Thus, there is a $z \in C(G)$ such that $d(z, G) = d(v_1, T)$, i.e., there is a size-preserving bijection $g: B(z, G) \rightarrow B(v_1, T)$.

If $m(v_1, T) < n$, then, by the minimality of n , every $D \in B(z, G)$ is properly isomorphic to $g(D)$, and thus G is isomorphic to T , which contradicts our assumption. Thus $m(v_1, T) \geq n$. Since all the branches of T at v_1 except $B(v_1, v_0, T)$ are properly contained in C and $e(C) = n$, we conclude that $m(v_1, T) = e(B(v_1, v_0, T))$. Clearly

$$(9) \quad n \leq e(B(v_1, v_0, T)) = k - e(C) + 1 = k - e(B) + 1 = e(B(u_1, u_0, G)),$$

and since all the branches of G at u_1 except $B(u_1, u_0, G)$ are properly contained in B , and $e(B) = n$, we conclude that

$$m(u_1, G) = e(B(u_1, u_0, G)) = m(v_1, T).$$

Since $d(z, G) = d(v_1, T)$, clearly $m(z, G) = m(v_1, T) = m(u_1, G)$. Thus, since G satisfies (iv), $d(u_1, G) = d(z, G) = d(v_1, T)$ and therefore there is a size-preserving bijection $h: B(u_1, G) \rightarrow B(v_1, T)$. Because of (9) we may assume that

$$h(B(u_1, u_0, G)) = B(v_1, v_0, T).$$

If $D \in B(u_1, G) \setminus B(u_1, u_0, G)$ then $e(D) < n$ and thus, by the minimality of n , D is properly isomorphic to $h(D)$. Therefore $UB_{u_0}(u_1, G)$ is isomorphic to $UB_{v_0}(v_1, T)$ by an isomorphism that carries u_1 onto v_1 . This isomorphism can clearly be extended to a proper isomorphism from B to C , which is impossible, since B and C , as chosen above, are not properly isomorphic. This completes Step IV and establishes Theorem 1. \square

Remark. The equivalence between conditions (i) and (iv) of Theorem 1 implies Theorem 2 of [1] and Theorems 3, 4, 5 and 6 of [2] as special cases.

References

- [1] N. Alon, *A note on the decomposition of trees into isomorphic subtrees*. Ars Combinatoria 12 (1981), 117-121.
- [2] Y. Caro and J. Schonheim, *Decomposition of trees into isomorphic subtrees*. Ars Combinatoria 9 (1980), 119-130.

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